a. For each entry [j,d,p] in M/Note $d=D_j^kp=p_j^k$) Set D_{ik}^k-d and $p_{ik}^k\leftarrow p$ b. For each destination, with an entry in M,

1. If event is a message M from neighbor k, {Called by PATH to process an event. }

Procedure NTU

{Invoked when the node comes up.} Procedure INIT-PATH

1. Initialize all tables.

2. Run PATH algorithm. End INIT-PATH

Algorithm PATH

Invoked when a message M is received from neighbor k, or an adjacent link to k has changed or when a node is

 Run NTU to update neighbor tables. initialized.}

Run MTU to update main tables.

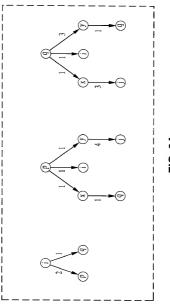
Add update entry $[j, D_i^* p_j^*]$ to the new message M. 4. Within finite amount of time, send message M to For each destination / marked as changed,

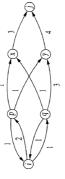
cach neighbor.

End PATH

2. If the event is an adjacent link-status change, update I_k^i and Remove existing links (n, j) in T_k' and add new link (m, j, d) to T_k' , where $d = D_{jk}' + D_{mk}'$ clear neighbor tables of k, if link is down. and $m = p_{jk}^i$.

End NTU





Preferred Neighbors

Destination p q x y

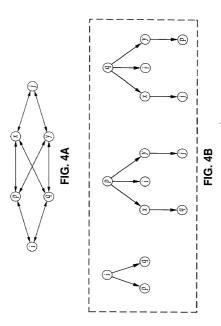
FIG. 3B

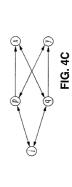
3 P

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7

Distance Pref. Nbr





Procedure MTU1. Clear link the Γ 2. For each node $j \neq i$ coursing in at least one T_i^i a. Find M IN $V = min(D_{j_i} + j_i)$ for V_{i_i} b. Let n be such that M IN $= (D_{j_i} + j_i)$. Thes are broken constitrently. Neighbor n is the preferred neighbor for destination j. For each link (j_i, v_i, d) in T_i^i Add link (j_i, v_i, d) to T3. Update T with each link l_i^i 4. Ann Dijstru's shortest path algorithm on T to find now $D_{j_i}^i$, and $p_{j_i}^i$ 5. For each electrication i if D_i or p_i changed from previous value, set changed and $cr_iport-i$ flags for j.

<u>.</u>

```
Procedure INIT-MPATH
      {Invoked when the node comes up.}
      1. Initialize tables and run MPATH.
End INIT-MPATH
Algorithm MPATH
      {Invoked when a message M is received from neighbor k,
     or an adjacent link to k has changed.}
      1. Run NTU to update neighbor tables.
     2. Run MTU to obtain new D_i^i and p_i^i
     3. If node is PASSIVE or node is ACTIVE ∧ last reply arrived,
           Reset goactive flag.
           For each destination j marked as report-it,
            a. FD_i^i \leftarrow min\{D_i^i RD_i^i\}
            b. If D_j > RD_j, Set goactive flag.
            c. RD_i^i \leftarrow D_i^i
            d. Add [j, RD_i^i, p_i^i] to message M'.
            e. Clear report-it flag for j.
        Otherwise, the node is ACTIVE and waiting for more replies,
         For each destination j marked as changed,
            f. FD_i^i \leftarrow min\{D_i^i, FD_i^i\}
      4. For each destination j marked as changed,
        a. Clear changed flag for j
        b. S_i^i \leftarrow \{k | D_{jk}^i \le FD_i^i\}
      For each neighbor k,
        a. M'' \leftarrow M'
        b. If event is query from k, Set reply flag in M".
        c. If goactive set, Set query flag in M".
        d.If M'' non-empty, send M'' to k.
      6. If goactive set, become ACTIVE, otherwise
        become PASSIVE.
END MPATH
```

FIG. 6